

# Divisor graphs have arbitrary order and size

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## Abstract

A divisor graph  $G$  is an ordered pair  $(V, E)$  where  $V \subset \mathbb{Z}$  and for all  $u \neq v \in V$ ,  $uv \in E$  if and only if  $u \mid v$  or  $v \mid u$ . A graph which is isomorphic to a divisor graph is also called a divisor graph. In this note, we will prove that for any  $n \geq 1$  and  $0 \leq m \leq \binom{n}{2}$  then there exists a divisor graph of order  $n$  and size  $m$ . We also present a simple proof of the characterization of divisor graphs which is due to Chartran, Muntean, Saenpholpant and Zhang.

## 1 Introduction

The notion of divisor graph was first introduced by Singh and Santhosh [2]. A divisor graph  $G(V)$  is an ordered pair  $(V, E)$  where  $V \subset \mathbb{Z}$  and for all  $u \neq v \in V$ ,  $uv \in E$  if and only if  $u \mid v$  or  $v \mid u$ . A graph which is isomorphic to a divisor graph is also called a divisor graph. The main result of this note is the following theorem.

**Theorem 1** *For any  $n \geq 1$  and  $0 \leq m \leq \binom{n}{2}$  then there exists a divisor graph of order  $n$  and size  $m$ .*

To prove Theorem 1, we need the following characterization of divisor graphs is due to Chartran, Muntean, Saenpholpant and Zhang [1].

**Theorem 2** *A graph  $G$  is divisor graph if and only if there is an orientation  $D$  of  $G$  such that if  $(x, y), (y, z)$  are edges of  $D$  then so is  $(x, z)$ .*

The proof of Theorem 2 in [1] is by induction on the order of the graph. For the completeness of this note, we will present a simple (and direct) proof of this theorem in Section 3. From Theorem 2, we introduce the definition of a divisor digraph which will be usefull in the proof of Theorem 1.

**Definition 1** *A digraph  $G$  is a divisor digraph if and only if  $(x, y), (y, z)$  are edges of  $G$  then so is  $(x, z)$ .*

It is clear that if  $G$  is a divisor digraph then the graph obtained by ignoring the direction of edges of  $G$  is a divisor graph.

## 2 Proof of Theorem 1

Suppose that  $G = (V, E)$  is a graph with vertex set  $V = \{v_1, \dots, v_n\}$  and size  $m$ . The degree of a vertex  $v_i$  is the number of edges of  $G$  incident with  $v_i$ . Let  $d_1 \geq d_2 \geq \dots \geq d_n$  be the vertex degrees in non-increasing order and let  $f_G = |\{i : d_i \geq i\}|$ . Let  $e_i$  be the number of vertices with degree at least  $i$ , that is  $e_i = |\{j : d_j \geq i\}|$ . Then we have  $e_1 \geq e_2 \geq \dots \geq e_n$ . Moreover, we also have

$$d_i = |\{j : e_j \geq i\}|. \quad (1)$$

We have the following lemmas.

**Lemma 1** *Let  $n - 1 = d_1 \geq \dots \geq d_n \geq 1$  be a sequence of natural numbers and let  $e_i = |\{j : d_j \geq i\}|$ . Suppose that*

$$\sum_{i=1}^n d_i = 2m, \quad (2)$$

and

$$\sum_{i=1}^t d_i = \sum_{i=1}^t (e_i - 1) \quad (3)$$

for all  $1 \leq t \leq f = |\{i : d_i \geq i\}|$ . Then there exists a divisor graph of order  $n$  and size  $m$ .

**Proof** We construct a digraph  $G$  with vertex set  $v_1, \dots, v_n$  as follows. For  $1 \leq i \leq f$ , then  $(v_i, v_j)$  is an edge of  $G$  for  $i + 1 \leq j \leq d_i + 1$ . We first show that  $G$  is a graph of size  $m$ . It suffices to show that  $\deg(v_i) = d_i$  for all  $i$  (where  $\deg(v_i)$  is the number of incident edges of vertex  $v_i$ , regardless of their directions). We have three cases.

1. Suppose that  $1 \leq i \leq f$ . Then it is clear from the construction that  $\deg(v_i) = d_i$ .
2. Suppose that  $i = f + 1$ . Then  $d_i = d_{f+1} \leq f$  (by the definition of  $f$ ). From (3), we have  $d_j = e_j - 1$  for  $1 \leq j \leq f$ . If  $d_{f+1} < f$  then  $e_f \leq f$ , or

$$f \leq d_f = e_f - 1 < f,$$

which is a contradiction. Hence  $d_{f+1} = f$ . For  $1 \leq j \leq f$  then  $d_j + 1 \geq d_f + 1 \geq f + 1 = i$ . So  $(v_j, v_i)$  is an edge of  $G$  and  $\deg(v_i) = f$ . Thus,  $\deg(v_i) = d_i$ .

3. Suppose that  $i > f + 1$ . Then  $(v_j, v_i)$  is an edge of  $G$  if and only if  $1 \leq j \leq f$  and  $j + 1 \leq i \leq d_j + 1$ . For  $j > f$  we have  $d_j \leq d_{f+1} < f + 1 \leq i$ . This implies that

$$\deg(v_i) = |\{1 \leq j \leq f : d_j \geq i - 1\}| = |\{j : d_j \geq i - 1\}| = e_{i-1}$$

for all  $i > f + 1$ . From (1), we have  $d_i = e_i - 1$  for  $1 \leq i \leq f$ . For  $i > f + 1$  then  $d_j \geq i - 1$  or  $e_j \geq i$  only if  $j \leq f$  (since  $d_{f+1} = f < i - 1$ ). Hence

$$\begin{aligned} e_{i-1} &= |\{j : d_j \geq i - 1\}| \\ &= |\{1 \leq j \leq f : d_j \geq i - 1\}| \\ &= |\{1 \leq j \leq f : e_j - 1 \geq i - 1\}| \\ &= |\{j : e_j \geq i\}| = d_i. \end{aligned}$$

Thus,  $\deg(v_i) = e_{i-1} = d_i$  for all  $i > f + 1$ .

Therefore, we have  $\deg(v_i) = d_i$  for  $1 \leq i \leq n$ . This implies that  $G$  has order  $n$  and size  $m$ .

Now, we will show that  $G$  is a divisor digraph. Suppose that  $(v_i, v_j)$  and  $(v_j, v_k)$  are two edges of  $G$ . Then from the above construction,  $1 \leq i, j \leq d$  and  $k \leq d_j + 1 \leq d_i + 1$ . Thus  $(v_i, v_k)$  is also an edge of  $G$ . This implies that  $G$  is a divisor digraph. Let  $H$  be the graph obtained from  $G$  by ignoring the direction of edges of  $G$ . Then  $H$  is a divisor graph of order  $n$  and size  $m$ . This concludes the proof of the lemma.  $\square$

**Lemma 2** *Let  $n - 1 = d_1 \geq \dots \geq d_n \geq 1$  be a sequence of natural numbers and set  $e_i = |\{j : d_j \geq i\}|$ . Suppose that*

$$\sum_{i=1}^n d_i = 2m < n(n-1), \quad (4)$$

and

$$\sum_{i=1}^t d_i = \sum_{i=1}^t (e_i - 1) \quad (5)$$

for all  $1 \leq t \leq f = |\{i : d_i \geq i\}|$ . Then there exists a sequence  $n - 1 = d'_1 \geq \dots \geq d'_n \geq 1$  of natural numbers such that

$$\sum_{i=1}^n d'_i = 2(m+1), \quad (6)$$

and

$$\sum_{i=1}^t d'_i = \sum_{i=1}^t (e'_i - 1) \quad (7)$$

for all  $1 \leq t \leq f' = |\{i : d'_i \geq i\}|$ , where  $e'_i = |\{j : d'_j \geq i\}|$  for  $1 \leq i \leq n$ .

**Proof** If  $f \geq n - 1$  then  $n - 1 \geq d_1 \geq \dots \geq d_{n-1} \geq d_f \geq f = n - 1$ . Hence  $d_1 = \dots = d_{n-1} = n - 1$ , and

$$\sum_{i=1}^{n-1} e_i = n(n-1).$$

We have  $e_i \leq n$  for  $1 \leq i \leq n - 1$ , so  $e_1 = \dots = e_{n-1} = n$ . Hence  $d_i \geq n - 1$  for  $1 \leq i \leq n$  or

$$\sum_{i=1}^n d_i = n(n-1),$$

which is a contradiction. Thus,  $f < n - 1$ . Let  $g$  be the smallest index such that  $d_i < n - 1$ . Then we have  $2 \leq g \leq f + 1$  (since  $d_{f+1} \leq f < n - 1$ ). We have two cases.

1. Suppose that  $2 \leq g \leq f$ . Set  $h = d_g + 2$ ,  $d'_g = d_g + 1$ ,  $d'_h = d_h + 1$  and  $d'_i = d_i$  for  $i \neq g, h$ . We have (6) holds since

$$\sum_{i=1}^n d'_i = 2 + \sum_{i=1}^n d_i = 2(m+1).$$

Recall that from the proof of Lemma 1, we have

$$d_i = \begin{cases} e_i - 1 & \text{if } i \leq f \\ f & \text{if } i = f + 1 \\ e_{i-1} & \text{if } i > f + 1. \end{cases} \quad (8)$$

Since  $d_g = h - 2 \geq g$  and  $g \leq f$ , we have  $e_g = h - 1$ . This implies that  $d_{h-1} \geq g$  and  $d_h < g$ . Besides,  $d_1 = \dots = d_{g-1} = n - 1$  so  $d_h \geq d_n = e_{n-1} \geq g - 1$ . Hence  $d_h = g - 1$ ,  $d'_g = h - 1$  and  $d'_h = g$ . Therefore, we have  $e'_g = e_g + 1$ ,  $e'_h = e_h + 1$  and  $e'_i = e_i$  for  $i \neq g, h$ . We have  $d_h = g - 1 \leq f - 1$  so  $h > f + 1$ . Hence  $d'_{f+1} = d_{f+1} < f + 1$  or  $f' \leq f$ . And we have (7) holds for  $1 \leq t \leq f'$  since (5) holds for  $1 \leq t \leq f' \leq f$ .

2. Suppose that  $g = f + 1$ . Then from (8) we have  $d_g = f = g - 1$ . Set  $h = f + 2$ ,  $d'_g = d_g + 1$ ,  $d'_h = d_h + 1$  and  $d'_i = d_i$  for  $i \neq g, h$ . Then it is clear that (6) holds. Besides, we have  $d_1 = \dots = d_{g-1} = n - 1$  so

$$g - 1 = d_g \geq \dots \geq d_n = e_{n-1} \geq g - 1.$$

Hence  $d_g = \dots = d_n = g - 1$ . We have  $d'_g = d'_h = g < h$ , so  $f' = f + 1$ ,  $e'_g = e_g + 2 = g + 1$  and  $e'_i = e_i$  for  $i \neq g$ . From (5), we have

$$\sum_{i=1}^t d'_i = \sum_{i=1}^t (e'_i - 1)$$

for  $1 \leq t \leq f$ . We only need to check for  $t = f' (= f + 1 = g)$ . We have

$$\sum_{i=1}^g d'_i = g + \sum_{i=1}^f d'_i = e'_g - 1 + \sum_{i=1}^f (e'_i - 1) = \sum_{i=1}^g (e'_i - 1).$$

Thus, (7) holds for  $1 \leq t \leq f'$ .

This concludes the proof of the lemma.  $\square$

We are now ready to prove Theorem 1. From Lemma 1, we start with the sequence  $(n - 1, 1, \dots, 1)$  to obtain a divisor graph of size  $n - 1$ . Then apply Lemma 1 and Lemma 2 inductively to obtain divisor graphs of order  $n, \dots, \binom{n}{2}$ . To construct a divisor graph of order  $n$  and size  $m$  with  $m < n - 1$ , we choose a vertex and join it with  $m$  other vertices. Thus, there exists a divisor graph of order  $n$  and size  $m$  for any  $n$  and  $0 \leq m \leq \binom{n}{2}$ . This concludes the proof of the theorem.

**Remark 1** An interesting and open question is to find necessary and sufficient conditions for a non-increasing sequence  $n - 1 \geq d_1 \geq \dots \geq d_n \geq 1$  such that there exists a divisor graphs with degree sequence  $(d_1, \dots, d_n)$ .

### 3 Proof of Theorem 2

Suppose that  $G$  is a divisor graph. Then there exists a set  $V$  of positive integer such that  $G \simeq G(V)$ . We give an orientation on each edge  $(i, j)$  of  $G$  as follows

$$(i, j) \in E(G), i \rightarrow j \text{ if and only if } i \mid j.$$

Suppose that  $(x, y), (y, z)$  are edges of  $D$ . Then  $x \mid y$  and  $y \mid z$ . Hence  $x \mid z$  and  $(x, z)$  is an edge of  $G$ .

Now suppose that there exists an orientation  $D$  of  $G$  such that if  $(x, y), (y, z)$  are edges of  $D$  then so is  $(x, z)$ . We will show that  $G$  is a divisor graph. We will give an explicit labelling for  $G$ . We start with any vertex of  $G$  and label it by  $\{a_1\}$  (a list of one symbol). Suppose that we have labelled  $k$  vertices of  $G$  and we have used  $a_1, \dots, a_l$  symbols (each vertex is labelled by a list of symbols and we will update this list in each step). We choose any unlabelled vertex, says  $v$ . Consider two sets

$$\begin{aligned} D_I(v) &= \{u \in V(G) \mid (u, v) \in E(D)\}, \\ D_O(v) &= \{u \in V(G) \mid (v, u) \in E(D)\}. \end{aligned}$$

We label  $v$  by  $L_v = \{a_{l+1}\}$ . For each  $u \in D_I(v)$  and  $u$  was labelled by a list  $L_u$  then we add this list into the list  $L_v$  to have a new list  $L_v$  for  $v$ . And for each  $u \in D_O(v)$  which was labelled by a list  $L_u$  then we add the new list  $L_v$  into  $L_u$  to have a new list for  $u$ . For each updated vertex  $u$ , we consider the set  $D_O(u)$ . For each  $w \in D_O(u)$  which was labelled by a list  $L_w$ , we add the new list  $L_u$  into the list  $L_w$  to have a new list for  $w$ . We keep doing until we have no vertex to update or we come back to some vertex which we met along the way. But in the latter case, we have a sequence of vertices, says  $w_1, \dots, w_t$  such that  $w_1 \rightarrow w_2 \rightarrow \dots \rightarrow w_t \rightarrow w_1$  in  $D$ . This implies that

$$w_1 \mid w_2, w_2 \mid w_3, \dots, w_t \mid w_1.$$

Hence  $w_1 = \dots = w_t$ , which is a contradiction. Thus the process must be stopped. We repeat the process until all the vertices of  $G$  have been labelled by lists of symbols. Suppose that we have used  $r$  symbols  $a_1, \dots, a_r$ . We choose  $r$  distinct primes  $p_1, \dots, p_r$  and for each vertex  $v \in V(G)$  which is labelled by a list  $\{a_{i_1}, \dots, a_{i_v}\} \subseteq \{a_1, \dots, a_r\}$  then we label it by the number

$$n(v) = p_{i_1} \dots p_{i_v}.$$

From the construction above, if  $(u, v)$  is an edge of  $G$  then either  $L_u \subset L_v$  or  $L_v \subset L_u$ . This implies that either  $u \mid v$  or  $v \mid u$ . Hence  $G$  is a divisor graph. This concludes the proof of the theorem.

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## References

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